Efficiency in percentage

Throughput in ops/s

Overhead in percentage

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ZHT: A Zero-hop Distributed Hash Table for High-end Computing Systems

Tongli Li1, Antonio Perez de Tejada1, Kevin Brandstatter1, Zhao Zhang2, Ioan Raicu1,3

1Department of Computer Science, Illinois Institute of Technology
2Department of Computer Science, University of Chicago
3Mathematics and Computer Science Division, Argonne National Laboratory

Abstract

One critical component of future file systems for high-end computing is meta-data management. This work presents ZHT, a zero-hop distributed hash table, which has been tuned for the requirements of HEC systems. ZHT aims to be a building block for future distributed file systems to implement distributed metadata management. The goals are delivering availability, fault tolerance, high throughput, and low latencies. ZHT has some important properties, such as being light-weight, fault tolerant using replication and persistence. We have evaluated ZHT’s performance under a variety of systems, ranging from a Linux cluster to an IBM BlueGene/P supercomputer. We scaled ZHT up to 16K processes and achieved 4M operations/sec throughput. Latencies have scaled similarly well, with sub-milliseconds latencies at 4K-core scales. We compared ZHT against other systems and found it offers superior performance for the features and portability it supports.

Hardware

- IBM Blue Gene/P supercomputer
- 1024 nodes
- 2GB RAM/node
- 4096 cores in total

Software

- OS: ZepTOS
- Batch execution system: Cobalt
- Persistent hash table: NoVoHT

Data serialization: Google protocol buffers

Experiment

Each client creates a long list of key-value pairs, here we set the length of key is 16 byte and length of value is 132 bytes. Clients firstly sequentially send all these key-value pairs ZHT API. ZHT will decide which server to send. Secondly clients send the same list of keys as lookup parameter to servers, finally send remove request with the same list of keys.

Evaluation

- More than 4M operations/sec aggregated throughput with 16K ZHT instances on 1024 nodes. As low as 0.78ms latency on 1024 node scale.
- UDP shows a better scalability and reliability. TCP can be as fast as UDP when using connection caching.
- The performance differences among three basic operations (insert, lookup and remove) are very small.

ZHT uses a direct 0-hop algorithm and that the majority of the overhead comes from network communication, it is not expected that the time per operation to increase significantly with larger scales.

Related work and acknowledgements


B. Fingersh, “Distributed caching with Memcached.” Linux Journal, 2004(124), 2004


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Future work

Broadcast primitive: transmit the key/value pair over the edges of the spanning tree with the goal to distribute the key/value pair to all the caches.

Data Indexing: When dealing with massive data collections, one challenge is indexing the material to support re-use and analysis.

Distributed Metadata Management: FusionFS will use ZHT to implement the distributed metadata management.