

HTCaaS: Leveraging Distributed Supercomputing Infrastructures for Large-Scale Scientific Computing

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Introduction

HTCaaS: High-Throughput Computing as a Service

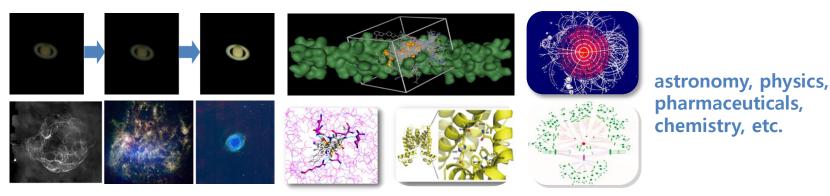
CEvaluation

Conclusions & Discussions



From HTC to Many-Task Computing (MTC) [MTAGS'08]

- > A very large number of tasks (millions or even billions)
- Relatively short per task execution times (sec to min)
- Data intensive tasks (i.e., tens of MB of I/O per second)
- A large variance of task execution times (i.e., ranging from hundreds of milliseconds to hours)
- Communication-intensive, however, not based on message passing interface (such as MPI) but through files





Middleware Systems for HTC/MTC applications

Ease of Use

✓ Minimize user overhead for handling jobs and resources

> Efficient Task Dispatching

✓ The overhead of task dispatching should be low enough

> Adaptiveness

✓ adjust acquired resources according to changing load

Fairness

✓ ensure fairness among multiple users submitting various numbers of tasks

> Reliability

✓ Failed or suspended tasks should be automatically resubmitted and managed

> Resource Integration

✓ effectively integrate as many computing resources as possible

on Clouds,

Introduction

Our Approach

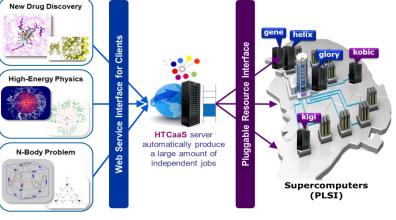
> High-Throughput Computing As a Service

✓ *Meta-Job* based automatic job split & submission

- e.g., parameter sweeps or N-body calculations
- ✓ Agent-based *multi-level scheduling*
- ✓ User-level Scheduling and *Dynamic Fairness*
- ✓ *Pluggable interface* to heterogeneous computing resources
- ✓ Supporting many *client interfaces*
 - Native WS-interface, Java API
 - Easy-to-use client tools (CLI/GUI/Web portal)

HTCaaS is currently running as a pilot service on top of PLSI

 ✓ supporting a number of scientific applications from pharmaceutical domain and high-energy physics









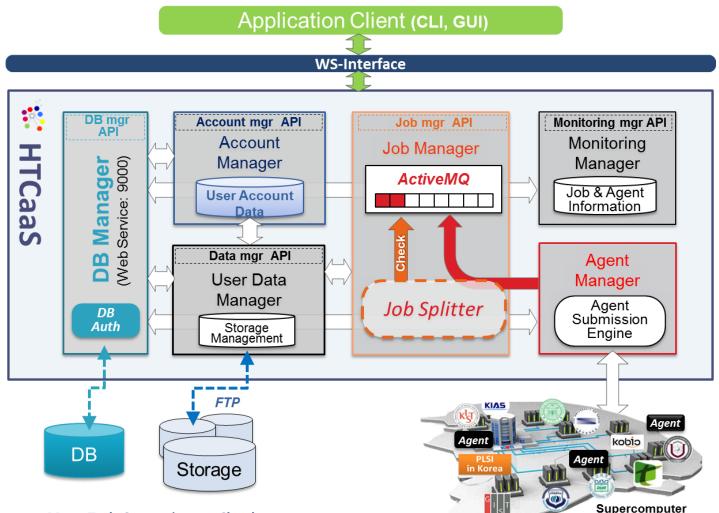
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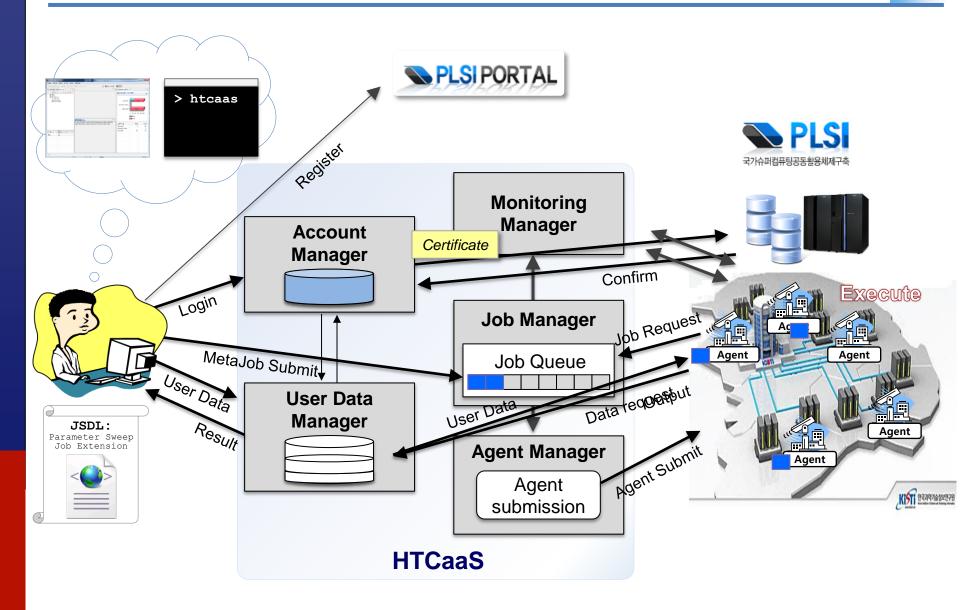
HTCaaS: High-Throughput Computing as a Service

Contraction

Conclusions & Discussions

System Architecture





Solution Job Queues and Agents Management

- > maintains separate job queues and agents per user
 - ✓ reducing complexities of *accounting* and *scheduling*
 - track and meter the usage of PLSI computing resources per user
 - carefully calibrating the number of agents per user can address the problem of *fair resource sharing* among multiple users
 - ✓ Each agent actively pulls the tasks from its dedicated job queue which corresponds to a specific user
 - if there are no more tasks to be processed, it automatically *releases* the acquired resources and exits

Monitoring and Fault-tolerance

- > periodically checks the status of agents and tasks
 - ✓ If some of agents or tasks fail, the Monitoring Manager informs the Agent Manager (or the Job Manager) to resubmit the failed agents (tasks) and manage them (addressing Reliability)

Our Content of Scheduling and Dynamic Fairness

> Dynamic fair resource sharing algorithm

- ✓ divides all available computing resources *fairly* across all *demanding* users in the system (when the system is heavily loaded) and exploits *dynamic adjustment* of acquired resources as free computing resources become available (as the overall system becomes lightly loaded)
- ✓ Resource Allotment Function

 $RA(U) = min\left(NumTasks(U), \frac{AvailableCores}{\sum_{p \in DU} Weight(p)} * Weight(U)\right)$

- Weight(p) represents the weight of a user p
 - » can consider many different factors such as the number of tasks submitted by the user p, task running time, priority, etc.
 - » Weighted Fairness

✓ addressing Fairness and Adaptiveness

	Timeline & Events		Computing Resources		RA	Job Queues
t ₀	All 100 cores are free	100 free cores				
t ₁	User A arrives and submits 200 tasks	User A's 100 tasks		RA(A)=100	A: 100 tasks(DU)	
t ₂	User B arrives and submits 50 tasks			RA(A)=50 RA(B)=50		
t ₃	User A release 50 cores					A: 50 tasks(DU)
t ₄ t ₅ t ₆	User B acquires 50 cores User C arrives and submits 20 tasks User B's tasks finish	User A's 50 tasks	User B's 50 tasks		RA(A)=50 RA(C)=20	B: 0 tasks(SU) A: 50 tasks(DU)
t ₇	User C acquires 20 cores	User A's 50 tasks	C's 20 tasks	30 free cores		C: 20 tasks(DU) A: 20 tasks(DU)
t ₈	User A acquires additional 30 cores	User A's 80 tasks		C's 20 tasks		C: 0 tasks(SU)

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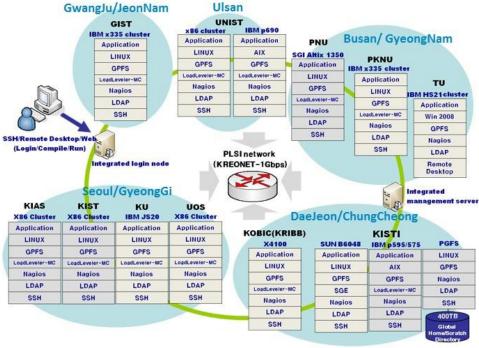
Conclusions & Discussions



PLSI (Partnership & Leadership for the nationwide Supercomputing Infrastructure)

provide researchers with an integrated view of geographically distributed supercomputing infrastructures to solve complex and demanding scientific problems

consisting of **multiple** organizations connected via a dedicated **1Gbps** network [**100TFLops** of computing power, 1,115 nodes with 8,508 CPU cores]





CPLSI provides a common software stack

- Accounting (based on LDAP), Monitoring (via Nagios), Global scheduling (based on LoadLeveler) and a Global shared storage system (based on GPFS)
 - ✓ utilize the LoadLeveler as a job submission system to available computing nodes in the PLSI
 - configured as a *multi-cluster* environment
 - ✓ exploits a total 400TB of global home/scratch directories mounted at *every* computing node as a shared storage system for input/output data and executables

ORG	SYSTEM	PROCESSOR	NETWORK	OS	CORES	MEM(GB)	GFLOPS
KIAS	helix (x86 cluster)	AMD Opteron 2GHz	1GbE	CentOS 6.2	128	8	1,024
KIAS	gene (x86 cluster)	AMD Opteron 2GHz	1 GbE	CentOS 6.2	128	8	1,024
KOBIC	kobic (SUN X4100)	AMD Opteron 2.1GHz	1 GbE	CentOS 5.4	184	4	1,545.6
KISTI	glory (SUN x2100)	AMD Opteron 1.8GHz	1GbE	CentOS 5.4	514	2	3,700.8

Table 1: PLSI Computing Resources leveraged by HTCaaS



OMicro-Benchmark Experiments

- simulating a large number of short-running tasks (sleep 10) and relatively long-running tasks (sleep 100)
- > glory cluster in KISTI (300 cores)
- Performance Metrics
 - ✓ Makespan (time to complete a bag of tasks)
 - ✓ Efficiency (comparison with ideal parallelism)

 $Efficiency(NT) = \frac{\frac{NT*PTE}{NCU}}{Makespan(NT)} * 100$

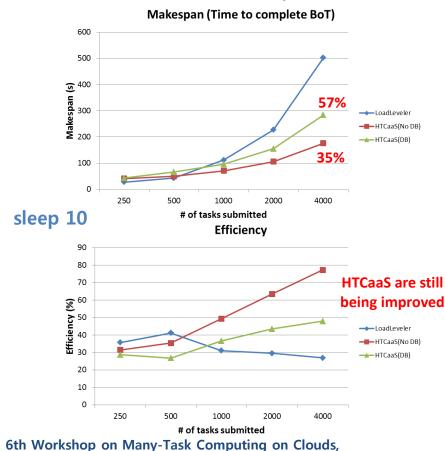
✓ Comparison Models

- HTCaaS with DB Manager connections (HTCaaS(DB))
- HTCaaS without DB Manager interactions (HTCaaS(No DB))
- LoadLeveler+GPFS (LoadLeveler)

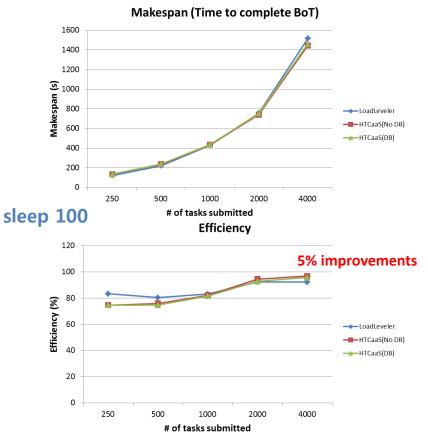


Micro-Benchmark Experiments

- > For short running tasks, HTCaaS clearly outperforms LoadLeveler
- For relatively long running tasks, overheads of task dispatching can be effectively countervailed



Grids, and Supercomputers (MTAGS'13)



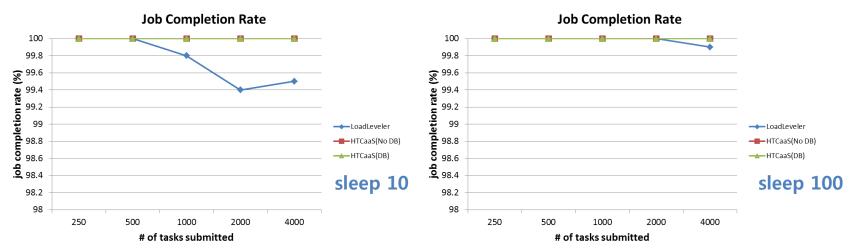


S Micro-Benchmark Experiments

Job holding problem during the course of LoadLeveler dispatching

 \checkmark due to multiple simultaneous I/O operations on the GPFS

- HTCaaS can effectively leverage the local storage of the computing resource
 - ✓ User Data Manager manages overall input/output data staging
 - support data-intensive HTC/MTC applications where typical size of a single input data is relatively small (from hundreds of KBs to MBs)





Protein Docking Experiments

- > Autodock, a suite of automated docking tools
 - ✓ perform the docking of *ligands* to a set of target *proteins* to discover new drugs for several serious diseases such as SARS or Malaria
- Experimental setup
 - ✓ Clusters: glory, kobic, gene and helix
 - ✓ Four different users are sequentially arriving at our system and submit various numbers of tasks (ligands) with an average inter-arrival time of 10 minutes
 - from 1000 down to 100 for the single cluster
 - from 2000 down to 250 for the multi-cluster
 - ✓ Comparison Models
 - HTCaaS with dynamic fair resource sharing algorithm dynamic fairness (**DF**)
 - Simple resource partitioning algorithm *strict* fairness (Simple)



Protein Docking Experiments

- HTCaaS can *reduce* the performance gap among users with various numbers of tasks
- HTCaaS can *seamlessly* integrate multiple geographically distributed clusters
 - ✓ fully utilize available computing resources as they become available, while Simple inevitably wastes idle computing resources

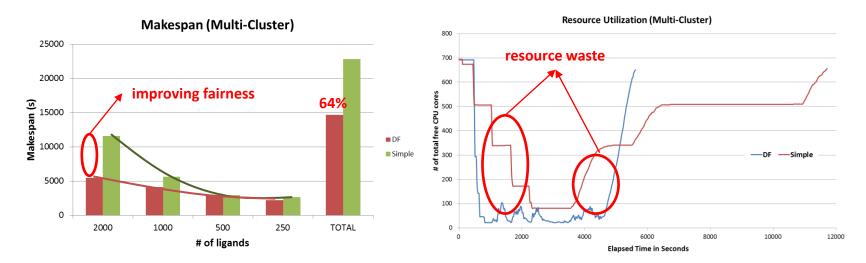


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Conclusion



HTCaaS to provide researchers with ease of exploring large-scale and complex HTC/MTC problems by integrating distributed national supercomputers

- Employing the concept of Meta-Job (with easy-to-use client tools) and a fault-tolerant agent-based multi-level scheduling mechanism (Ease of Use, Efficient Task Dispatching and Reliability)
- Employing dynamic fair resource sharing mechanism (Fairness and Adaptiveness)
- HTCaaS effectively leverages local disks of geographically distributed computing resources

✓ support data-intensive HTC/MTC applications

HTCaaS can seamlessly utilize all of available computing resources without resource wastage (Resource Integration)

Discussion



\$Future Work

- Supporting more complex workloads consisting of HTC and HPC tasks
- improving the scalability of HTCaaS, and applying job profiling technique to realize the weighted form of fairness
- How we can more efficiently support *data-intensive* MTC applications on top of *geographically distributed* computing environments such as PLSI?
 - ✓ Shared storage system such as GPFS will be performance bottleneck
 - ✓ Data caching may work but it depends on the workload characteristics



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